

Quaternary Cryptographic Bent Functions and their Binary projection

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The aim of this work

m-variables quaternary functions(F, m)
 $F : GR(4, m) \rightarrow \mathbb{Z}_4 \quad \mathbb{Z}_4 = \{0, 1, 2, 3\}$

n – variables boolean functions(f, n)
 $f : \mathbb{F}_2^n \rightarrow \mathbb{F}_2 \quad (\mathbb{F}_2 = \{0, 1\})$

The quaternary approach (CONSTRUCTION):

- GALOIS rings $R = GR(4, m)$.
- CRYPTOGRAPHIC properties .
- Characterization of a family of quaternary CRYPTOGRAPHIC functions.

From \mathbb{Z}_4 to \mathbb{F}_2 (APPLICATION):

- The binary projection.
- Drived boolean cryptographic functions.

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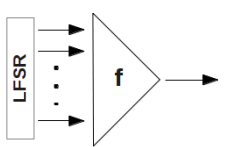
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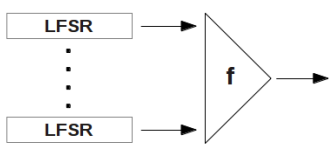
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- 1 Boolean and Quaternary functions
- 2 Galois Rings
- 3 The construction
- 4 Derived boolean functions
- 5 Complete example of construction

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Let $f : \mathbb{F}_2^n \mapsto \mathbb{F}_2$

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- Support : $\text{supp}(f) = \{u \in \mathbb{F}_2^n \mid f(u) \neq 0\}$
- Weight : $w_H(f) = |\text{supp}(f)|$
- HAMMING metric : $d_H(f, g) = w_H(f \oplus g)_{(\oplus = + \text{mod}(2))}$

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CRYPTOGRAPHIC PROPERTIES.

$$\text{Walsh transform : } W_f(u) = \sum_{v \in \mathbb{F}_2^n} (-1)^{u \cdot v} (-1)^{f(v)}, u \in \mathbb{F}_2^n$$

$$\begin{array}{l} \text{Balancedness} : w_H(f) = 2^{n-1} \iff W_f(0) = 0 \\ \text{Nonlinearity} : nl_2(f) = \min_{g \text{ affine}} d_H(f, g) \iff nl_2(f) = 2^{n-1} - \frac{1}{2} \max_{u \in \mathbb{F}_2^n} |W_f(u)| \end{array}$$

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$$f \text{ is bent} \iff \forall u \in \mathbb{F}_2^n, |W_f(u)| = 2^{\frac{n}{2}}.$$

Maximal nonlinearity $(2^{n-1} - 2^{\frac{n}{2}-1})$ for n even but (not balanced).

Quaternary functions

Let $F, G \in \mathcal{F}(\mathbb{Z}_4^m, \mathbb{Z}_4) : \mathbb{Z}_4^m \mapsto \mathbb{Z}_4$

The ring of integers modulus 4

$$\mathbb{Z}_4 = \mathbb{Z}/4\mathbb{Z} = \{0, 1, 2, 3\}$$

The group of 4th root of unity in \mathbb{C}

$$\mathbb{U}_4 = \{\pm 1, \pm i\} \quad i^2 = -1$$

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- Relative support : $\text{supp}_j(F) = \{u \in \mathbb{Z}_4^m \mid F(u) = j\}_{0 \leq j \leq 3}$
- Relative cardinal : $\eta_j(F) = |\text{supp}_j(F)|$

LEE METRIC

$z \in \mathbb{Z}_4$	0	1	2	3
$w_L(z)$	0	1	2	1

- LEE Weight : $w_L(F) = \eta_1(F) + 2\eta_2(F) + \eta_3(F)$
- LEE Distance : $d_L(F, G) = w_L(F - G) \pmod{4}$

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WALSH TRANSFORM.

$$W_F(u) = \sum_{v \in \mathbb{Z}_4^n} i^{u \cdot v} j^{F(v)}, \quad u \in \mathbb{Z}_4^n$$

$$W_F^2(u) = \sum_{v \in \mathbb{Z}_4^m} i^{u \cdot v} (-1)^{F(v)}, \quad u \in \mathbb{Z}_4^n$$

CRYPTOGRAPHIC PROPERTIES.

Let $F \in \mathcal{F}(\mathbb{Z}_4^m, \mathbb{Z}_4)$

The function F is **balanced**

$$\begin{aligned} \iff \eta_j(F) &= 4^{m-1} \quad \forall j \in \{0, 1, 2, 3\} \\ \iff W_F(0) &= W_F^2(0) = 0 \end{aligned}$$

The **nonlinearity** of F :

$$\begin{aligned} nl_4^L(F) &= \min_{G \text{ affine}} d_L(F, G) \\ &= 4^m - \max_{a \in \mathbb{Z}_4^m, b \in \mathbb{Z}_4} \left\{ \text{Re}(i^b W_F(a)) \right\} \end{aligned}$$

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F is bent $\iff \forall x \in \mathbb{Z}_4^m, |W_F(x)| = 2^m$

The nonlinearity of a **bent** function F is $nl_4^L(F) = 4^m - 2^m$.

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GALOIS Rings

THE MULTIPLICATIVE REPRESENTATION AND CYCLOTOMIC CLASSES

$$R \simeq \mathbb{Z}_4[x]/(g(x)) \quad (\text{b_poly of deg } m) \simeq \mathbb{Z}_4[\beta] \quad (2^m - 1^{\text{th}} \text{ root of unity}) \simeq \mathbb{Z}_4^m$$

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$$GR(4, m) = R^* / T \cup D \cup T$$

$$\forall z \in R, \quad z = \begin{cases} \beta^j & \text{if } z \in T^* \\ 2\beta^j & \text{if } z \in D^* \\ \beta^i + 2\beta^j & \text{if } z \in R^* / T \\ 0 & \text{if } z = 0 \end{cases} \quad \text{(Multiplicative representation)} \quad 0 \leq j, i \leq 2^m - 2$$

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The 2^m -CYCLOTOMIC CLASSES $(C_j)_{0 \leq j \leq 2^m - 1}$ of order $2^m - 1$ of R^* are defined by :

$$\begin{cases} C_j = \{\beta^l(1 + 2\beta^j), 0 \leq l \leq 2^m - 2\} \\ C_{2^m - 1} = \{\beta^l, 0 \leq l \leq 2^m - 2\} \end{cases}$$

$$R = \bigcup_{j=0}^{2^m - 1} C_j \cup D$$

▶ Return

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The m -variables quaternary function

DEFINITION OF F_k :

$$\forall k \in \{0, 1, \dots, 2^m - 2\}$$

We define the m -variables quaternary function F_k as follow :

$$\begin{array}{lcl}
 F_k & : & R \rightarrow \mathbb{Z}_4 \\
 & & a + 2b \rightarrow F_k(a + 2b) = h_k(\beta^k(1 + 2ba^{2^m-2}))
 \end{array}$$

where, $h_k : \mathfrak{C}_k \rightarrow \mathbb{Z}_4$ and $\mathfrak{C}_k = \{\beta^k\} \cup \{\beta^k(1 + 2\beta^j), 0 \leq j \leq 2^m - 2\}$

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CHARACTERISATION OF F_k :

$$F_k(x) = \begin{cases} h_k(\beta^k(1 + 2\beta^j)) & \text{if } x \in C_j, 0 \leq j \leq 2^m - 2. \\ h_k(\beta^k) & \text{if } x \in C_{2^m-1} \cup D. \end{cases}$$

► Go to

Conditions on the intern function h_k

Relation between bentness of F_k and his intern function h_k

$$\forall k \in \{0, 1, \dots, 2^m - 2\}$$

The q-ary function F_k is bent if

$$\forall x \in \mathcal{T} \begin{cases} \sum_{v \in \mathcal{C}_k} i^{h_k(v)} = 0 & (1) \\ \left| \sum_{v \in \mathcal{T}} i^{h_k(\beta^k(1+2v))+3\text{Tr}(v \oplus x)} \right| = 2^{\frac{m}{2}} & (2) \end{cases}$$

with,

$$\forall a, b \in \mathcal{T}, \quad a \oplus b = a + b + 2(ab)^{2^{m-1}} \in \mathcal{T}$$

$$\forall z \in \mathcal{R}, \quad \text{Tr}(z) = \sum_{l=0}^{m-1} \sigma^l(z) \in \mathbb{Z}_4$$

Algebraic duality

Let note E^* the dual of \mathcal{T} , and $L(x, y) : \mathcal{T} \times \mathcal{T} \rightarrow 2\mathbb{F}_2$ a bilinear, symmetric and nondegenerate function.

$$\text{if, } \forall x \in \mathcal{T} \begin{cases} l_1(x) = L(b_1, x) \\ l_2(x) = L(b_2, x) \\ l_3(x) = L(b_3, x) \end{cases}$$

are three balanced functions
where $b_1 \neq b_2 \neq 0$ and,
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Then the orthogonal of $B = \{0, b_1, b_2, b_3\}$ is :

$$B^\perp = \{x \in \mathcal{T} \mid \forall l \in B^* : l(x) = 0\}, \quad B^* = \{l_1, l_2\}$$

And, if $\forall i \in \{1, 2, 3\} \quad b_i \notin B^\perp \implies |B^\perp| = 2^{m-2}$

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Example of splitting:

Let $L(x, y) = 2\text{Tr}(xy)$ and $b_1 \neq b_2 \in \mathcal{T}^*$ with $\text{Tr}(b_1)$ is even and $\text{Tr}(b_1 b_2)$ is odd Then :

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$$\mathcal{T} = B^\perp \oplus B = \cup_{i=0}^3 b_i \oplus B^\perp$$

CONSTRUCTION OF h_K

Let $\alpha = (\alpha_0, \alpha_1, \alpha_2, \alpha_3) \in \mathbb{Z}_4^4$ such that:

$$(S_1) \quad \begin{cases} |j^{\alpha_0} + j^{\alpha_1} + j^{\alpha_2} + j^{\alpha_3}| = 2 \\ |j^{\alpha_0} + j^{\alpha_1} - j^{\alpha_2} - j^{\alpha_3}| = 2 \\ |j^{\alpha_0} - j^{\alpha_1} + j^{\alpha_2} - j^{\alpha_3}| = 2 \\ |j^{\alpha_0} - j^{\alpha_1} - j^{\alpha_2} + j^{\alpha_3}| = 2 \end{cases}$$

and $b_1, b_2 \in \mathcal{T}^*$ defined in the splitting example.

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$$\forall 0 \leq j \neq i \leq 3, \quad \alpha_i + \text{Tr}(b_i) \neq \alpha_j + \text{Tr}(b_j) \quad (S_2)$$

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Then the intern function h_k defined by :

$$\forall j, 0 \leq j \leq 3, \quad \forall x \in B^\perp \oplus b_j, \quad h_k(\beta^k(1 + 2x)) = \alpha_j + \text{Tr}(b_j)$$

verifies the conditions (1) and (2) .

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Let α a solution of (S_2) then $\alpha + i \pmod 4$ and $\sigma^{13}(\alpha)$ are also solutions.

Solutions and models

We note a **model** \mathbf{a} the vector (a_0, a_1, a_2, a_3) defined by:

$$\forall 0 \leq i \neq j \leq 3, \quad a_j = \alpha_j + \text{Tr}(b_j) \quad \text{and} \quad a_i \neq a_j.$$

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The quaternary **Bent** function, $F_k : R = \cup_{j=0}^{2^m-1} C_j \cup D \rightarrow \mathbb{Z}_4$

$$\begin{cases} F_k(C_j) & = a_j \text{ if } \beta^j \in B^\perp \oplus b_j \\ F_k(D \cup \mathcal{T}^*) & = a_0 \end{cases}$$

outline

- 1 Boolean and Quaternary functions
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- 3 The construction
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- 5 Complete example of construction

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$$\forall z \in R: z = \sum_{j=0}^{m-1} z_j \beta^j \quad \begin{array}{l} z_j \in \mathbb{Z}_4 \\ \text{(Additive representation)} \\ \text{with, } \beta \text{ a root of the } b\text{-polynomial} \end{array}$$

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REDEFINITION OF F_k ($k \in \{0, 1, \dots, 2^m - 2\}$):

$$\begin{array}{l} \bar{F}_k : E^* \cup W \cup (\cup_{j=0}^{2^m-2} V_j) \rightarrow \mathbb{Z}_4 \\ \quad \quad \quad u + 2v \quad \quad \quad \rightarrow \bar{F}_k(u + 2v) = F_k(d^{-1}(u + 2v)) \end{array}$$

$$\bar{h}_k : d(\mathfrak{e}_k) \rightarrow \mathbb{Z}_4, \bar{h}_k(x) = h_k(d^{-1}(x))$$

The direct mapping between \mathbb{Z}_4^m and \mathbb{F}_2^{2m}

FROM \mathbb{Z}_4^m TO \mathbb{F}_2^{2m} :

$$\begin{array}{lcl} \varphi : & \mathbb{Z}_4^m & \rightarrow \mathbb{F}_2^{2m} \\ & u + 2v & \rightarrow \tilde{u} \parallel \tilde{v} \end{array}$$

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where $\tilde{\cdot}$ is the component mod 2 reduction and \parallel is the concatenation.

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Let $\psi_{i \in \mathbb{N}}$ any mapping from \mathbb{Z}_4 to \mathbb{F}_2 such that :

$$\sum_{x \in \mathbb{Z}_4} (-1)^{\psi_i(x)} = 0$$

Example: $\psi_i(2q+r) = q$ or $\psi_i(2q+r) = q+r \pmod{2}$.

The $2m$ -variables Derived boolean functions

DEFINITION OF f :

The $2m$ -variables boolean function f derived from the quaternary constructed function F_k and defined as :

$$\begin{array}{lcl}
 f & : & \mathbb{F}_2^{2m} \rightarrow \mathbb{F}_2 \\
 x & \mapsto & \psi_i(\bar{F}_k(\varphi^{-1}(x))) \text{ is bent.}
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$$f(\tilde{u}||\tilde{v}) = \begin{cases} \psi_i(\bar{h}_k(v_k(v_0 + 2v_j))) & \text{if } \varphi^{-1}(\tilde{u}||\tilde{v}) \in V_{j_{0 \leq j \leq 2^m-2}} \\ \psi_i(\bar{h}_k(v_k)) & \text{if } \varphi^{-1}(\tilde{u}||\tilde{v}) \in E^* \cup W \end{cases}$$

The $2m + 1$ -variables derived boolean functions

DEFINITION OF f :

The derived boolean function f defined as :

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 f : \mathbb{F}_2^{2m+1} &\rightarrow \mathbb{F}_2 \\
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has maximal nonlinearity equal to $4^m - 2^{m+1}$.

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\mathcal{T}^* E^*	1 { 1, 0, 0 }	β { 0, 1, 0 }	β^2 { 0, 0, 1 }	β^3 { 1, 3, 2 }	β^4 { 2, 3, 3 }	β^5 { 3, 3, 1 }	β^6 { 1, 2, 1 }
D^* W^*	2 { 2, 0, 0 }	2β { 0, 2, 0 }	$2\beta^2$ { 0, 0, 2 }	$2\beta^3$ { 2, 2, 0 }	$2\beta^4$ { 0, 2, 2 }	$2\beta^5$ { 2, 2, 2 }	$2\beta^6$ { 2, 0, 2 }
C_0	3	3β	$3\beta^2$	$3\beta^3$	$3\beta^4$	$3\beta^5$	$3\beta^6$
C_1	$1 + 2\beta$	$\beta + 2\beta^2$	$\beta^2 + 2\beta^3$	$\beta^3 + 2\beta^4$	$\beta^4 + 2\beta^5$	$\beta^5 + 2\beta^6$	$\beta^6 + 2$
C_2	$1 + 2\beta^2$	$\beta + 2\beta^3$	$\beta^2 + 2\beta^4$	$\beta^3 + 2\beta^5$	$\beta^4 + 2\beta^6$	$\beta^5 + 2$	$\beta^6 + 2\beta$
C_3	$1 + 2\beta^3$	$\beta + 2\beta^4$	$\beta^2 + 2\beta^5$	$\beta^3 + 2\beta^6$	$\beta^4 + 2$	$\beta^5 + 2\beta$	$\beta^6 + 2\beta^2$
C_4	$1 + 2\beta^4$	$\beta + 2\beta^5$	$\beta^2 + 2\beta^6$	$\beta^3 + 2$	$\beta^4 + 2\beta$	$\beta^5 + 2\beta^2$	$\beta^6 + 2\beta^3$
C_5	$1 + 2\beta^5$	$\beta + 2\beta^6$	$\beta^2 + 2$	$\beta^3 + 2\beta$	$\beta^4 + 2\beta^2$	$\beta^5 + 2\beta^3$	$\beta^6 + 2\beta^4$
C_6	$1 + 2\beta^6$	$\beta + 2$	$\beta^2 + 2\beta$	$\beta^3 + 2\beta^2$	$\beta^4 + 2\beta^3$	$\beta^5 + 2\beta^4$	$\beta^6 + 2\beta^5$

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C_0 V_0	3 {3, 0, 0}	3β {0, 3, 0}	$3\beta^2$ {0, 0, 3}	$3\beta^3$ {3, 1, 2}	$3\beta^4$ {2, 1, 1}	$3\beta^5$ {1, 1, 3}	$3\beta^6$ {3, 2, 3}
C_1 V_1	$1 + 2\beta$ {1, 2, 0}	$\beta + 2\beta^2$ {0, 1, 2}	$\beta^2 + 2\beta^3$ {2, 2, 1}	$\beta^3 + 2\beta^4$ {1, 1, 0}	$\beta^4 + 2\beta^5$ {0, 1, 1}	$\beta^5 + 2\beta^6$ {1, 3, 3}	$\beta^6 + 2$ {3, 2, 1}
C_2 V_2	$1 + 2\beta^2$ {1, 0, 2}	$\beta + 2\beta^3$ {2, 3, 0}	$\beta^2 + 2\beta^4$ {0, 2, 3}	$\beta^3 + 2\beta^5$ {3, 1, 0}	$\beta^4 + 2\beta^6$ {0, 3, 1}	$\beta^5 + 2$ {1, 3, 1}	$\beta^6 + 2\beta$ {1, 0, 1}
C_3 V_3	$1 + 2\beta^3$ {3, 2, 0}	$\beta + 2\beta^4$ {0, 3, 2}	$\beta^2 + 2\beta^5$ {2, 2, 3}	$\beta^3 + 2\beta^6$ {3, 3, 0}	$\beta^4 + 2$ {0, 3, 3}	$\beta^5 + 2\beta$ {3, 1, 1}	$\beta^6 + 2\beta^2$ {1, 2, 3}
C_4 V_4	$1 + 2\beta^4$ {1, 2, 2}	$\beta + 2\beta^5$ {2, 3, 2}	$\beta^2 + 2\beta^6$ {2, 0, 3}	$\beta^3 + 2$ {3, 3, 2}	$\beta^4 + 2\beta$ {2, 1, 3}	$\beta^5 + 2\beta^2$ {3, 3, 3}	$\beta^6 + 2\beta^3$ {3, 0, 1}
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Choice of b_1 and b_2

$b_i \in \mathcal{T}$	1	β	β^2	β^3	β^4	β^5	β^6	0
$\text{Tr}(b_i)$	3	2	2	1	2	1	1	0

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b_2									
1		0	β^3	β^6	β	β^5	β^4	β^2	1
β		β^3	0	β^4	1	β^2	β^6	β^5	β
β^2		β^6	β^4	0	β^5	β	β^3	1	β^2
β^3		β	1	β^5	0	β^6	β^2	β^4	β^3
β^4		β^5	β^2	β	β^6	0	1	β^3	β^4
β^5		β^4	β^6	β^3	β^2	1	0	β	β^5
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\oplus	b_1	1	β	β^2	β^3	β^4	β^5	β^6	0
b_2									
1		0	β^3	β^6	β	β^5	β^4	β^2	1
β		β^3	0	β^4	1	β^2	β^6	β^5	β
β^2		β^6	β^4	0	β^5	β	β^3	1	β^2
β^3		β	1	β^5	0	β^6	β^2	β^4	β^3
β^4		β^5	β^2	β	β^6	0	1	β^3	β^4
β^5		β^4	β^6	β^3	β^2	1	0	β	β^5
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β^3		β	1	β^5	0	β^6	β^2	β^4	β^3
β^4		β^5	β^2	β	β^6	0	1	β^3	β^4
β^5		β^4	β^6	β^3	β^2	1	0	β	β^5
β^6		β^2	β^5	1	β^4	β^3	β	0	β^6
0		1	β	β^2	β^3	β^4	β^5	β^6	0

We have $\mathcal{T} = \cup_{i=0}^3 B^\perp \oplus b_i$ and $B^\perp = \{x \in \mathcal{T}, 2Tr(xb_1) = 2Tr(xb_2) = 0\}$

$$\begin{aligned}
 B^\perp \oplus b_0 &= \{0, \beta^6\} \\
 B^\perp \oplus b_1 &= \{\beta^2, 1\} \\
 B^\perp \oplus b_2 &= \{\beta^3, \beta^4\} \\
 B^\perp \oplus b_3 &= \{\beta^5, \beta\}
 \end{aligned}$$

Then we can define explicitly our quaternary function like that:

$$\forall x \in R, F_k(x) = \begin{cases} 0 & \text{if } x \in UC_6 \cup D \cup T^* \\ 2 & \text{if } x \in UC_2 \cup C_0 \\ 1 & \text{if } x \in UC_3 \cup C_4 \\ 3 & \text{if } x \in UC_5 \cup C_1 \end{cases}$$

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Then \bar{F} is defined such that:

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$$\forall x \in \mathbb{Z}_4^m, f(\varphi(x)) = \begin{cases} 0 & \text{if } \varphi(x) \in \cup \varphi(V_6) \cup \varphi(W) \cup \varphi(E^*) \\ 1 & \text{if } \varphi(x) \in \cup \varphi(V_2) \cup \varphi(V_0) \\ 0 & \text{if } \varphi(x) \in \cup \varphi(V_3) \cup \varphi(V_4) \\ 1 & \text{if } \varphi(x) \in \cup \varphi(V_5) \cup \varphi(V_1) \end{cases}$$

Let $\psi_{\varepsilon \in \{0,1\}} : \mathbb{Z}_4 \rightarrow \mathbb{F}_2$, $\psi_{\varepsilon}(2q + r) = q * \varepsilon + \bar{\varepsilon} * r$ then:

The $2m + 1$ -Derived boolean function it's defined like that:

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The $2m + 1$ -Derived boolean function it's defined like that:

$$\forall x \in \mathbb{Z}_4^m, f(\varphi(x) || \varepsilon) = \begin{cases} 0 & \text{if } \varphi(x) || \varepsilon \in \cup \varphi(V_6)_{\varepsilon} \cup \varphi(W)_{\varepsilon} \cup \varphi(E^*)_{\varepsilon} \\ \varepsilon & \text{if } \varphi(x) || \varepsilon \in \cup \varphi(V_2)_{\varepsilon} \cup \varphi(V_0)_{\varepsilon} \\ \bar{\varepsilon} & \text{if } \varphi(x) || \varepsilon \in \cup \varphi(V_3)_{\varepsilon} \cup \varphi(V_4)_{\varepsilon} \\ \varepsilon + \bar{\varepsilon} & \text{if } \varphi(x) || \varepsilon \in \cup \varphi(V_5)_{\varepsilon} \cup \varphi(V_1)_{\varepsilon} \end{cases}$$

Conclusion

- THE QUATERNARY FUNCTION:

$$\begin{array}{lcl}
 F_k & : & R = GR(4, m) \rightarrow \mathbb{Z}_4 \\
 & & a + 2b \rightarrow F_k(a + 2b) = h_k(\beta^k(1 + 2ba^{2^m-2}))
 \end{array}$$

with $\forall k, 0 \leq k \leq 2^m - 2, h_k : \mathfrak{C}_k \rightarrow \mathbb{Z}_4$

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- THE BINARY PROJECTIONS:

$$f : \mathbb{F}_2^{2m} \rightarrow \mathbb{F}_2$$

$$x \rightarrow \psi_i(F_k(d^{-1}(\varphi^{-1}(x))))$$

$$f : \mathbb{F}_2^{2m+1} \rightarrow \mathbb{F}_2$$

$$x || \varepsilon \rightarrow \psi_\varepsilon(F_k(d^{-1}(\varphi^{-1}(x))))$$

with $R \xrightarrow{d} \mathbb{Z}_4^m \xrightarrow{\varphi} \mathbb{F}_2^{2m}$

and $\psi, \psi_\varepsilon : \mathbb{Z}_4 \xrightarrow{\text{balanced}} \mathbb{F}_2$